

# Convergence Rates in the Probabilistic Analysis of Algorithms

Ralph Neininger

Institute for Mathematics, Goethe University, 60054 Frankfurt a.M., Germany  
neininger@math.uni-frankfurt.de

Jasmin Straub

Institute for Mathematics, Goethe University, 60054 Frankfurt a.M., Germany  
jstraub@math.uni-frankfurt.de

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## Abstract

In this extended abstract a general framework is developed to bound rates of convergence for sequences of random variables as they mainly arise in the analysis of random trees and divide and conquer algorithms. The rates of convergence are bounded in the Zolotarev distances. Concrete examples from the analysis of algorithms and data structures are discussed as well as a few examples from other areas. They lead to convergence rates of polynomial and logarithmic order. A crucial role is played by a factor 3 in the exponent of these orders in cases where the normal distribution is the limit distribution.

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## 1 Introduction and notation

In this extended abstract we consider a general recurrence for (probability) distributions which covers many instances of complexity measures of divide and conquer algorithms and parameters of random search trees. We consider a sequence  $(Y_n)_{n \geq 0}$  of  $d$ -dimensional random vectors satisfying the distributional recursion

$$Y_n \stackrel{d}{=} \sum_{r=1}^K A_r(n) Y_{I_r^{(n)}}^{(r)} + b_n, \quad n \geq n_0, \quad (1)$$

where  $(A_1(n), \dots, A_K(n), b_n, I^{(n)})$ ,  $(Y_n^{(1)})_{n \geq 0}, \dots, (Y_n^{(K)})_{n \geq 0}$  are independent, the coefficients  $A_1(n), \dots, A_K(n)$  are random  $(d \times d)$ -matrices,  $b_n$  is a  $d$ -dimensional random vector,  $I^{(n)} = (I_1^{(n)}, \dots, I_K^{(n)})$  is a random vector in  $\{0, \dots, n\}^K$ ,  $n_0 \geq 1$  and  $(Y_n^{(r)})_{n \geq 0} \stackrel{d}{=} (Y_n)_{n \geq 0}$  for  $r = 1, \dots, K$ . Moreover,  $K \geq 1$  is a fixed integer, but extensions to  $K$  being random and depending on  $n$  are possible.

This is the framework of [14] where some general convergence results are shown for appropriate normalizations of the  $Y_n$ . The content of the present extended abstract is to also study the rates of convergence in such limit theorems.

We define the normalized sequence  $(X_n)_{n \geq 0}$  by

$$X_n := C_n^{-1/2}(Y_n - M_n), \quad n \geq 0,$$

where  $M_n$  is a  $d$ -dimensional vector and  $C_n$  a positive definite  $(d \times d)$ -matrix. Essentially, we choose  $M_n$  as the mean and  $C_n$  as the covariance matrix of  $Y_n$  if they exist or as the



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41 leading order terms in expansions of these moments as  $n \rightarrow \infty$ . The normalized quantities  
42 satisfy the following modified recursion:

$$43 \quad X_n \stackrel{d}{=} \sum_{r=1}^K A_r^{(n)} X_{I_r^{(n)}}^{(r)} + b^{(n)}, \quad n \geq n_0, \quad (2)$$

44  
45 with

$$46 \quad A_r^{(n)} := C_n^{-1/2} A_r(n) C_{I_r^{(n)}}^{1/2}, \quad b^{(n)} := C_n^{-1/2} \left( b_n - M_n + \sum_{r=1}^K A_r(n) M_{I_r^{(n)}} \right) \quad (3)$$

47  
48 and independence relations as in (1).

49 In the context of the contraction method the aim is to establish transfer theorems of the  
50 following form: After verifying the assumptions of appropriate convergence of the coefficients  
51  $A_r^{(n)} \rightarrow A_r^*$ ,  $b^{(n)} \rightarrow b^*$  then convergence in distribution of random vectors  $(X_n)$  to a limit  $X$   
52 is implied. The limit distribution  $\mathcal{L}(X)$  is identified by a fixed-point equation obtained from  
53 (2) by considering formally  $n \rightarrow \infty$ :

$$54 \quad X \stackrel{d}{=} \sum_{r=1}^K A_r^* X^{(r)} + b^*. \quad (4)$$

55  
56 Here  $(A_1^*, \dots, A_K^*, b^*)$ ,  $X^{(1)}, \dots, X^{(K)}$  are independent and  $X^{(r)} \stackrel{d}{=} X$  for  $r = 1, \dots, K$ .

57 The aim of the present extended abstract is to endow such general transfer theorems  
58 with bounds on the rates of convergence. As a distance measure between (probability)  
59 distributions we use the Zolotarev metric. For various of the applications we discuss, bounds  
60 on the rate of convergence have been derived one by one for more popular distance measures  
61 such as the Kolmogorov–Smirnov distance. The transfer theorems of the present paper are  
62 in terms of the smoother Zolotarev metrics. However, they are easy to apply and cover a  
63 broad range of applications at once. A crucial role is played by a factor 3 in the exponent of  
64 these orders in cases where the normal distribution is the limit distribution, see Remark 4.

65 In the rest of this section we fix some notation. Regarding norms of vectors and (random)  
66 matrices we denote for  $x \in \mathbb{R}^d$  by  $\|x\|$  its Euclidean norm and for a random vector  $X$  and  
67 some  $0 < p < \infty$ , we set  $\|X\|_p := \mathbb{E}[\|X\|^p]^{(1/p) \wedge 1}$ . Furthermore, for a  $(d \times d)$ -matrix  $A$ ,  
68  $\|A\|_{\text{op}} := \sup_{\|x\|=1} \|Ax\|$  denotes the spectral norm of  $A$  and for a random such  $A$  we define  
69  $\|A\|_p := \mathbb{E}[\|A\|_{\text{op}}^p]^{(1/p) \wedge 1}$  for a random square matrix and  $0 < p < \infty$ . Note that for a  
70 symmetric  $(d \times d)$ -matrix  $A$ , we have  $\|A\|_{\text{op}} = \max\{|\lambda| : \lambda \text{ eigenvalue of } A\}$ . By  $\text{Id}_d$  the  
71  $d$ -dimensional unit matrix is denoted. For multilinear forms the norm is defined similarly.

72 Furthermore we define by  $\mathcal{P}^d$  the space of probability distributions in  $\mathbb{R}^d$  (endowed with  
73 the Borel  $\sigma$ -field), by  $\mathcal{P}_s^d := \{\mathcal{L}(X) \in \mathcal{P}^d : \|X\|_s < \infty\}$  and for a vector  $m \in \mathbb{R}^d$ , and a  
74 symmetric positive semidefinite  $d \times d$  matrix  $C$  the spaces

$$75 \quad \mathcal{P}_s^d(m) := \{\mathcal{L}(X) \in \mathcal{P}_s^d : \mathbb{E}[X] = m\}, \quad s > 1, \quad (5)$$

$$76 \quad \mathcal{P}_s^d(m, C) := \{\mathcal{L}(X) \in \mathcal{P}_s^d : \mathbb{E}[X] = m, \text{Cov}(X) = C\}, \quad s > 2.$$

77  
78 We use the convention  $\mathcal{P}_s^d(m) := \mathcal{P}_s^d$  for  $s \leq 1$  and  $\mathcal{P}_s^d(m, C) := \mathcal{P}_s^d(m)$  for  $s \leq 2$ .

79 The Zolotarev metrics  $\zeta_s$ , [19], are defined for probability distributions  $\mathcal{L}(X), \mathcal{L}(Y) \in \mathcal{P}^d$   
80 by

$$81 \quad \zeta_s(X, Y) := \zeta_s(\mathcal{L}(X), \mathcal{L}(Y)) = \sup_{f \in \mathcal{F}_s} |E(f(X) - f(Y))| \quad (6)$$

82

83 where for  $s = m + \alpha$ ,  $0 < \alpha \leq 1$ ,  $m \in \mathbb{N}_0$ ,

$$84 \quad \mathcal{F}_s := \{f \in C^m(\mathbb{R}^d, \mathbb{R}) : \|f^{(m)}(x) - f^{(m)}(y)\| \leq \|x - y\|^\alpha\}.$$

86 Note that these distance measures may be infinite. Finite metrics are given by  $\zeta_s$  on  $\mathcal{P}_s^d$  for  
87  $0 \leq s \leq 1$ , by  $\zeta_s$  on  $\mathcal{P}_s^d(m)$  for  $1 < s \leq 2$ , and by  $\zeta_s$  on  $\mathcal{P}_s^d(m, C)$  for  $2 < s \leq 3$ , cf. (5).

## 88 2 Results

89 We return to the situation outlined in the introduction, where we have normalized  $(Y_n)_{n \geq 0}$   
90 in the following way:

$$91 \quad X_n := C_n^{-1/2}(Y_n - M_n), \quad n \geq 0, \quad (7)$$

93 where  $M_n$  is a  $d$ -dimensional random vector and  $C_n$  a positive definite  $(d \times d)$ -matrix. As  
94 recalled in Section 1, for  $s > 1$ , we may fix the mean and covariance matrix of the scaled  
95 quantities to guarantee the finiteness of the  $\zeta_s$ -metric. Therefore, we choose  $M_n = \mathbb{E}[Y_n]$   
96 for  $n \geq 0$  and  $s > 1$ . For  $s > 2$ , we additionally have to control the covariances of  $X_n$ . We  
97 assume that there exists an  $n_1 \geq 0$  such that  $\text{Cov}(Y_n)$  is positive definite for  $n \geq n_1$  and  
98 choose  $C_n = \text{Cov}(Y_n)$  for  $n \geq n_1$  and  $C_n = \text{Id}_d$  for  $n < n_1$ . For  $s \leq 2$ , we just assume that  
99  $C_n$  is positive definite and set  $n_1 = 0$  in this case.

100 The normalized quantities satisfy the modified recursion

$$101 \quad X_n \stackrel{d}{=} \sum_{r=1}^K A_r^{(n)} X_{I_r^{(n)}}^{(r)} + b^{(n)}, \quad n \geq n_0,$$

103 with  $A_r^{(n)}$  and  $b^{(n)}$  given in (3). The following theorem discusses a general framework to  
104 bound rates of convergence for the sequence  $(X_n)_{n \geq 0}$ . For the proof, we need some technical  
105 conditions which guarantee that the sizes  $I_r^{(n)}$  of the subproblems grow with  $n$ . More precisely,  
106 we will assume that there exists some monotonically decreasing sequence  $R(n) > 0$  with  
107  $R(n) \rightarrow 0$  such that

$$108 \quad \|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_s = O(R(n)), \quad n \rightarrow \infty, \quad (8)$$

110 for all  $\ell \in \mathbb{N}$  and  $r = 1, \dots, K$  and that

$$111 \quad \|\mathbf{1}_{\{I_r^{(n)} = n\}} A_r^{(n)}\|_s \rightarrow 0, \quad n \rightarrow \infty, \quad (9)$$

113 for all  $r = 1, \dots, K$ .

### 114 2.1 A general transfer theorem for rates of convergence

115 Our first result is a direct extension of the main Theorem 4.1 in [14], where we essentially  
116 only make all the estimates there explicit. The main result of the present extended abstract  
117 is contained in the subsequent subsection.

118 ► **Theorem 1.** *Let  $(X_n)_{n \geq 0}$  be  $s$ -integrable,  $0 < s \leq 3$ , and satisfy recurrence (7) with the  
119 choices for  $M_n$  and  $C_n$  specified there. We assume that there exist  $s$ -integrable  $A_1^*, \dots, A_K^*, b^*$   
120 and some monotonically decreasing sequence  $R(n) > 0$  with  $R(n) \rightarrow 0$  such that, as  $n \rightarrow \infty$ ,*

$$121 \quad \|b^{(n)} - b^*\|_s + \sum_{r=1}^K \|A_r^{(n)} - A_r^*\|_s = O(R(n)). \quad (10)$$

122

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123 If conditions (8) and (9) are satisfied and if

$$124 \quad \limsup_{n \rightarrow \infty} \mathbb{E} \sum_{r=1}^K \left( \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right) < 1, \quad (11)$$

125 then we have, as  $n \rightarrow \infty$ ,

$$127 \quad \zeta_s(X_n, X) = O(R(n)),$$

128 where  $\mathcal{L}(X)$  is given as the unique fixed point in  $\mathcal{P}_s^d(0, \text{Id}_d)$  of the equation

$$130 \quad X \stackrel{d}{=} \sum_{r=1}^K A_r^* X^{(r)} + b^*, \quad (12)$$

131 with  $(A_1^*, \dots, A_K^*, b^*), X^{(1)}, \dots, X^{(K)}$  independent and  $X^{(r)} \stackrel{d}{=} X$  for  $r = 1, \dots, K$ .

133 **► Remark 2.** In applications, the convergence rate of the coefficients (conditions (8) and  
134 (10)) is often faster than the convergence rate of the quantities  $X_n$ , see, e.g., Section 4.4.  
135 In these cases, it is often possible to perform the induction step in the proof of Theorem 1  
136 although condition (11) does not hold. To be more precise, we may assume

$$137 \quad \|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_s + \|b^{(n)} - b^*\|_s + \|A_r^{(n)} - A_r^*\|_s = O(\tilde{R}(n))$$

138 for every  $\ell \geq 0$ ,  $r = 1, \dots, K$  and  $n \rightarrow \infty$ . Then, instead of condition (11), it is sufficient to  
139 find some  $K > 0$  such that

$$140 \quad \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{n_1 \leq I_r^{(n)} < n\}} \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right] \leq 1 - p_n - \frac{\tilde{R}(n)}{KR(n)} \quad (13)$$

141 for all large  $n$  with  $p_n := \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)} = n\}} \|A_r^{(n)}\|_{\text{op}}^s \right]$ .

### 143 2.2 An improved transfer theorem for normal limit distributions

144 We now consider the special case where the sequence  $(X_n)_{n \geq 0}$  is 3-integrable and satisfies  
145 recursion (2) with  $(A_1^{(n)}, \dots, A_K^{(n)}, b^{(n)}) \xrightarrow{L_3} (A_1^*, \dots, A_K^*, b^*)$  for some 3-integrable coefficients  
146  $A_1^*, \dots, A_K^*, b^*$  with

$$147 \quad b^* = 0, \quad \sum_{r=1}^K A_r^* (A_r^*)^T = \text{Id}_d$$

148 almost surely. Corollary 3.4 in [14] implies that, if  $\mathbb{E}[\sum_{r=1}^K \|A_r^*\|_{\text{op}}^3] < 1$ , equation (12) has a  
149 unique solution in the space  $\mathcal{P}_3^d(0, \text{Id}_d)$ . Furthermore, e.g., using characteristic functions, it  
150 is easily checked that this unique solution is the standard normal distribution  $\mathcal{N}(0, \text{Id}_d)$ .

151 In this special case of normal limit laws, it is possible to derive a refined version of  
152 Theorem 1. Instead of the technical condition (8), we now need the weaker condition

$$153 \quad \|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_3^3 = O(R(n)), \quad n \rightarrow \infty, \quad (14)$$

154 for all  $\ell \in \mathbb{N}$  and  $r = 1, \dots, K$ . Moreover, condition (10) concerning the convergence rates of  
155 the coefficients can be weakened, which is formulated in the following theorem.  
156

157 ► **Theorem 3.** Let  $(X_n)_{n \geq 0}$  be given as in (7) and be 3-integrable. We assume that for some  
 158  $R(n) > 0$  monotonically decreasing with  $R(n) \rightarrow 0$  as  $n \rightarrow \infty$  we have

$$159 \quad \left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3 = O(R(n)), \quad (15)$$

160 and the technical conditions (9) and (14) being satisfied for  $s = 3$ . If

$$162 \quad \limsup_{n \rightarrow \infty} \mathbb{E} \sum_{r=1}^K \left( \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^3 \right) < 1, \quad (16)$$

163 then we have, as  $n \rightarrow \infty$ ,

$$165 \quad \zeta_3(X_n, \mathcal{N}(0, \text{Id}_d)) = O(R(n)).$$

166 **Proof.** (Sketch) We define an accompanying sequence  $(Z_n^*)_{n \geq 0}$  by

$$167 \quad Z_n^* := \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} N^{(r)} + b^{(n)}, \quad n \geq 0,$$

168 where  $(A_1^{(n)}, \dots, A_K^{(n)}, I^{(n)}, b^{(n)}, N^{(1)}, \dots, N^{(K)})$  are independent,  $\mathcal{L}(N^{(r)}) = \mathcal{N}(0, \text{Id}_d)$  for  
 169  $r = 1, \dots, K$  and  $T_n T_n^T = \text{Cov}(X_n)$  for  $n \geq 0$ . Hence,  $Z_n^*$  is  $L_3$ -integrable,  $\mathbb{E}[Z_n^*] = 0$  and  
 170  $\text{Cov}(Z_n^*) = \text{Id}_d$  for all  $n \geq n_1$ . By the triangle inequality, we have

$$171 \quad \zeta_3(X_n, \mathcal{N}(0, \text{Id}_d)) \leq \zeta_3(X_n, Z_n^*) + \zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)).$$

172 Then, the assertion follows inductively if one has shown the bound  $\zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) =$   
 173  $O(R(n))$ : Using the convolution property of the multidimensional normal distribution, we  
 174 obtain the representation

$$175 \quad Z_n^* = \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} N^{(r)} + b^{(n)} \stackrel{d}{=} G_n N + b^{(n)}, \quad (17)$$

177 where  $G_n G_n^T = \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} T_{I_r^{(n)}}^T (A_r^{(n)})^T$ ,  $\mathcal{L}(N) = \mathcal{N}(0, \text{Id}_d)$  and  $N$  is independent of  
 178  $(G_n, b^{(n)})$ . As  $\text{Cov}(Z_n^*) = \text{Id}_d$  for all  $n \geq n_1$ , we have  $\mathbb{E}[G_n G_n^T + b^{(n)} (b^{(n)})^T] = \text{Id}_d$  for  $n \geq n_1$ .  
 179 Furthermore, we have  $\|b^{(n)}\|_3^3 = O(R(n))$  and

$$\begin{aligned} 180 \quad \|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} &= \left\| \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} T_{I_r^{(n)}}^T (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \\ 181 \quad &= O \left( \left\| \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)} < n_1\}} A_r^{(n)} (T_{I_r^{(n)}} T_{I_r^{(n)}}^T - \text{Id}_d) (A_r^{(n)})^T \right\|_{3/2}^{3/2} \right. \\ 182 \quad &\quad \left. + \left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \right) \\ 183 \quad &= O \left( \sum_{r=1}^K \left\| \mathbf{1}_{\{I_r^{(n)} < n_1\}} A_r^{(n)} \right\|_3^3 + \left\| \sum_{r=1}^K A_r^{(n)} (A_r^{(n)})^T - \text{Id}_d \right\|_{3/2}^{3/2} \right) \\ 184 \quad &= O(R(n)). \end{aligned}$$

186 Thus, the following Lemma 5 implies  $\zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = O(R(n))$ . Lemma 5 is the main  
 187 part of the present proof. ◀

188 ▶ Remark 4. Theorem 3, when applicable, often improves over Theorem 1 by a factor 3 in  
 189 the exponent, see Remark 9 for an example. This is caused by the additional exponents in  
 190 (15) in comparison to (10).

191 ▶ Lemma 5. Let  $(Z_n^*)_{n \geq 0}$  be a sequence of  $d$ -dimensional random vectors satisfying  $Z_n^* \stackrel{d}{=} G_n N + b^{(n)}$ , where  $G_n$  is a random  $(d \times d)$ -matrix,  $b^{(n)}$  a centered random vector with  
 192  $\mathbb{E}[G_n G_n^T + b^{(n)}(b^{(n)})^T] = \text{Id}_d$  and  $N \sim \mathcal{N}(0, \text{Id}_d)$  independent of  $(G_n, b^{(n)})$ . Furthermore, we  
 193 assume that, as  $n \rightarrow \infty$ ,

$$195 \quad \|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3 = O(R(n))$$

196 for appropriate  $R(n)$ . Then, we have, as  $n \rightarrow \infty$ ,

$$197 \quad \zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = O(R(n)).$$

198 The proof of Lemma 5 builds upon ideas of [15].

### 199 3 Expansions of moments

200 In applications to problems arising in theoretical computer science, where the recurrence  
 201 (1) is explicitly given one usually has no direct means to identify the orders of the terms  
 202  $\|b^{(n)} - b^*\|_s$  and  $\|A_r^{(n)} - A_r^*\|_s$ . This is due to the fact that the mean vector  $M_n$  and the  
 203 covariance matrix  $C_n$ , for the cases  $1 < s \leq 2$  and  $2 < s \leq 3$  respectively, which are used  
 204 for the normalization (7) are typically not exactly known or too involved to be amenable  
 205 to explicit calculations. As a substitute one usually has asymptotic expansions of these  
 206 sequences as  $n \rightarrow \infty$ .

207 In the present section we assume the dimension to be  $d = 1$  and  $A_r(n) = 1$  for all  
 208  $r = 1, \dots, K$  and provide tools to apply the general Theorems 1 and 3 on the basis of  
 209 expansions of the mean and variance. We assume that

$$210 \quad \mathbb{E}[X_n] = \mu(n) = f(n) + O(e(n)), \quad \text{Var}(X_n) = \sigma^2(n) = g(n) + O(h(n)), \quad (18)$$

212 with  $e(n) = o(f(n))$  and  $h(n) = o(g(n))$ . To connect Theorems 1 and 3 to recurrences with  
 213 known expansions we use the following notion.

214 ▶ Definition 6. A sequence  $(a(n))_{n \geq 0}$  of non-negative numbers is called essentially non-  
 215 decreasing if there exists a  $c > 0$  such that  $a(m) \leq ca(n)$  for all  $0 \leq m < n$ .

216 The scaling introduced in (7) with the special choices  $A_r(n) = 1$  for all  $r = 1, \dots, K$  leads to  
 217 the scaled recurrence for  $(X_n)$  given in (2) with

$$218 \quad A_r^{(n)} = \frac{\sigma(I_r^{(n)})}{\sigma(n)}, \quad b^{(n)} = \frac{1}{\sigma(n)} \left( b_n - \mu(n) + \sum_{r=1}^K \mu(I_r^{(n)}) \right). \quad (19)$$

220 Additionally, we consider the corresponding quantities

$$221 \quad \bar{A}_r^{(n)} = \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)}, \quad \bar{b}^{(n)} = \frac{1}{g^{1/2}(n)} \left( b_n - f(n) + \sum_{r=1}^K f(I_r^{(n)}) \right). \quad (20)$$

223 Then we have:

224 ► **Lemma 7.** With  $A_r^{(n)}$ ,  $b^{(n)}$  given in (19),  $\bar{A}_r^{(n)}$ ,  $\bar{b}^{(n)}$  given in (20), and the expansions for  
 225  $\mu(n)$ ,  $\sigma^2(n)$  given in (18) the following holds.

226 If the sequence  $h/g^{1/2}$  is essentially non-decreasing then

$$227 \quad \|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)}\right). \quad (21)$$

228  
 229 If the sequence  $h$  is essentially non-decreasing then

$$230 \quad \left\| \sum_{r=1}^K (A_r^{(n)})^2 - 1 \right\|_s \leq \left\| \sum_{r=1}^K (\bar{A}_r^{(n)})^2 - 1 \right\|_s + O\left(\frac{h(n)}{g(n)}\right). \quad (22)$$

231  
 232 If the sequence  $e$  is essentially non-decreasing then

$$233 \quad \|b^{(n)} - b^*\|_s \leq \|\bar{b}^{(n)} - b^*\|_s + O\left(\frac{h(n)}{g(n)} + \frac{e(n)}{g^{1/2}(n)}\right). \quad (23)$$

234  
 235 If the sequence  $g/h$  is essentially non-decreasing and

$$236 \quad T(n) := \mathbb{E} \sum_{r=1}^K \frac{g^{s/2-1}(I_r^{(n)})h(I_r^{(n)})R(I_r^{(n)})}{g^{s/2}(n)R(n)}$$

237 then we have

$$238 \quad \mathbb{E} \sum_{r=1}^K \frac{\sigma^s(I_r^{(n)})R(I_r^{(n)})}{\sigma^s(n)R(n)} \leq \mathbb{E} \sum_{r=1}^K \frac{g^{s/2}(I_r^{(n)})R(I_r^{(n)})}{g^{s/2}(n)R(n)} + O(T(n)). \quad (24)$$

240 **Proof.** We show (21), the other bounds can be shown similarly. Note that  $\sigma^2(n) = g(n) +$   
 241  $O(h(n))$  implies  $\sigma(n) = g^{1/2}(n) + O(h(n)/g^{1/2}(n))$  and that for any essentially non-decreasing  
 242 sequence  $(a(n))_{n \geq 0}$  we have  $\|a(I_r^{(n)})\|_\infty = O(a(n))$ . Since  $h/g^{1/2}$  is essentially non-decreasing  
 243 we obtain

$$244 \quad A_r^{(n)} = \frac{\sigma(I_r^{(n)})}{\sigma(n)} = \frac{g^{1/2}(I_r^{(n)}) + O(h(I_r^{(n)})/g^{1/2}(I_r^{(n)}))}{\sigma(n)}$$

$$245 \quad = \frac{g^{1/2}(I_r^{(n)}) + O(h(n)/g^{1/2}(n))}{g^{1/2}(n)} \cdot \frac{g^{1/2}(n)}{\sigma(n)}$$

$$246 \quad = \left( \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} + O\left(\frac{h(n)}{g(n)}\right) \right) \left( 1 + O\left(\frac{h(n)}{g(n)}\right) \right)$$

$$247 \quad = \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} + O\left(\frac{h(n)}{g(n)} \left( 1 + \frac{g^{1/2}(I_r^{(n)})}{g^{1/2}(n)} \right)\right).$$

248  
 249 Hence, we obtain

$$250 \quad \|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)} \left( 1 + \|\bar{A}_r^{(n)}\|_s \right)\right).$$

251  
 252 Since  $\bar{A}_r^{(n)} \rightarrow A_r^*$  in  $L_s$  we have  $\|\bar{A}_r^{(n)}\|_s = O(1)$ , hence

$$253 \quad \|A_r^{(n)} - A_r^*\|_s \leq \|\bar{A}_r^{(n)} - A_r^*\|_s + O\left(\frac{h(n)}{g(n)}\right),$$

254  
 255 which is bound (21). ◀

256 Note that in applications the terms on the right hand side in the estimates (21)–(24) can  
 257 easily be bound when expansions as in (18) with explicit functions  $e, f, g, h$  are available.

258 **4 Applications**

259 We start by deriving a known result to illustrate in detail how to apply our framework of the  
 260 previous sections.

261 **4.1 Quicksort: Key comparisons**

262 The number of key comparisons  $Y_n$  needed by the Quicksort algorithm to sort  $n$  randomly  
 263 permuted (distinct) numbers satisfies the distributional recursion

264 
$$Y_n \stackrel{d}{=} Y_{I_n} + Y'_{n-1-I_n} + n - 1, \quad n \geq 1, \quad (25)$$

265 where  $Y_0 := 0$  and  $(Y_k)_{k=0, \dots, n-1}, (Y'_k)_{k=0, \dots, n-1}, I_n$  are independent,  $I_n$  is uniformly distrib-  
 266 uted on  $\{0, \dots, n-1\}$ , and  $Y_k \stackrel{d}{=} Y'_k, k \geq 0$ . Hence, equation (25) is covered by our general  
 267 recurrence (1). For the expectation and variance of  $Y_n$  exact expressions are known which  
 268 imply the asymptotic expansions

269 
$$\mathbb{E}Y_n = 2n \log(n) + (2\gamma - 4)n + O(\log n), \quad (26)$$

270 
$$\text{Var}(Y_n) = \sigma^2 n^2 - 2n \log(n) + O(n), \quad (27)$$

272 where  $\gamma$  denotes Euler's constant and  $\sigma := \sqrt{7 - 2\pi^2/3} > 0$ . We introduce the normalized  
 273 quantities  $X_0 := X_1 := X_2 := 0$  and

274 
$$X_n := \frac{Y_n - \mathbb{E}Y_n}{\sqrt{\text{Var}(Y_n)}}, \quad n \geq 3. \quad (28)$$

275 To apply Theorem 1 we need to find an  $0 < s \leq 3$  and a sequence  $(R(n))$  with (10) and (11).  
 276 Note that the  $Y_n$  are bounded, thus  $L_s$ -integrable for any  $s > 0$ . To bound the  $L_s$ -norms  
 277 appearing in (10) we use Lemma 7 and choose

278 
$$f(n) = 2n \log(n) + (2\gamma - 4)n, \quad e(n) = \log n,$$

279 
$$g(n) = \sigma^2 n^2, \quad h(n) = n \log n.$$

281 With these functions we obtain for the quantities defined in (20) that

282 
$$\bar{A}_1^{(n)} = \frac{I_n}{n}, \quad \bar{A}_2^{(n)} = \frac{n-1-I_n}{n},$$

283 
$$\bar{b}^{(n)} = \frac{1}{\sigma} \left( 2 \frac{I_n}{n} \log \frac{I_n}{n} + 2 \frac{n-1-I_n}{n} \log \frac{n-1-I_n}{n} + \frac{n-1}{n} + O\left(\frac{\log n}{n}\right) \right)$$

285 With the embedding  $I_n = \lfloor nU \rfloor$  with  $U$  uniformly distributed over the unit interval  $[0, 1]$  we  
 286 have

287 
$$A_1^* = U, \quad A_2^* = 1 - U, \quad b^* = \frac{1}{\sigma} (2U \log(U) + 2(1 - U) \log(1 - U) + 1) =: \frac{1}{\sigma} \varphi(U).$$

289 The limit theorem  $X_n \rightarrow X$  has been derived by different methods by Régnier [16] and  
 290 Rösler [17]. Rösler [17] also found that the scaled limit  $Y := \sigma X$  satisfies the distributional  
 291 fixed-point equation

292 
$$Y \stackrel{d}{=} UY + (1 - U)Y' + \varphi(U). \quad (29)$$

293 Lower and upper bounds for the rate of convergence in  $X_n \rightarrow X$  have been studied for  
 294 various metrics in Fill and Janson [6] and Neininger and Rüschemdorf [13].



295 Now, we apply the framework of the present paper: For  $r = 1, 2$  and any  $s \geq 1$  we find  
 296 that

$$297 \quad \|\bar{A}_r^{(n)} - A_r^*\|_s = O\left(\frac{1}{n}\right).$$

298  
 299 Using Proposition 3.2 of Rösler [17] we obtain

$$300 \quad \|\bar{b}_n - b^*\|_s = O\left(\frac{\log n}{n}\right).$$

301  
 302 Moreover, we have

$$303 \quad \frac{h(n)}{g(n)} = O(R(n)) \quad \text{and} \quad \frac{e(n)}{g^{1/2}(n)} = O(R(n)) \quad \text{with} \quad R(n) := \frac{\log n}{n},$$

304  
 305 thus Lemma 7 implies that condition (10) is satisfied for our choice of the sequence  $R$ . To  
 306 verify condition (11) by use of (24) we obtain that for  $T(n)$  given in Lemma 7 we find  
 307  $T(n) = O(\log(n)/n) \rightarrow 0$  and that

$$308 \quad \mathbb{E} \sum_{r=1}^2 \frac{g^{s/2}(I_r^{(n)})R(I_r^{(n)})}{g^{s/2}(n)R(n)} = \mathbb{E} \sum_{r=1}^2 \left(\frac{I_r^{(n)}}{n}\right)^{s-1} \frac{\log I_r^{(n)}}{\log n}.$$

309  
 310 Note that the latter expression has a limes superior of less than 1 if and only if  $s > 2$ . Hence,  
 311 Theorem 1 is applicable for  $s > 2$  and yields that

$$312 \quad \zeta_s(X_n, X) = O\left(\frac{\log n}{n}\right), \quad \text{for} \quad 2 < s \leq 3. \tag{30}$$

313  
 314 The bound (30) had previously been shown for  $s = 3$  in [13], where also the optimality of  
 315 the order was shown, i.e., that  $\zeta_3(X_n, X) = \Theta(\log(n)/n)$ .

316 In the full paper version we also discuss bounds on rates of convergence for various cost  
 317 measures of the related Quickselect algorithms under various models for the rank to be  
 318 selected.

## 319 4.2 Size of $m$ -ary search trees

320 The size of  $m$ -ary search trees satisfies the recurrence (1) with  $K = m \geq 3$ ,  $A_1(n) = \dots =$   
 321  $A_m(n) = 1$ ,  $n_0 = m$ ,  $b_n = 1$ , i.e., we have

$$322 \quad Y_n \stackrel{d}{=} \sum_{r=1}^m Y_{I_r^{(n)}}^{(r)} + 1, \quad n \geq m.$$

323  
 324 For a representation of  $I^{(n)}$  we define for independent, identically unif[0, 1] distributed random  
 325 variables  $U_1, \dots, U_{m-1}$  their spacings in  $[0, 1]$  by  $S_1 = U_{(1)}, S_2 = U_{(2)} - U_{(1)}, \dots, S_m :=$   
 326  $1 - U_{(m-1)}$ , where  $U_{(1)}, \dots, U_{(m-1)}$  denote the order statistics of  $U_1, \dots, U_{m-1}$ . Then  $I^{(n)}$   
 327 has the mixed multinomial distribution:

$$328 \quad I^{(n)} \stackrel{d}{=} M(n - m + 1, S_1, \dots, S_m).$$

329 By this we mean that given  $(S_1, \dots, S_m) = (s_1, \dots, s_m)$  we have that  $I^{(n)}$  is multinomial  
 330  $M(n - m + 1, s_1, \dots, s_m)$  distributed. Expectations, variances and limit laws for  $Y_n$  have  
 331 been studied, see[12, 4]. We have

$$332 \quad \mathbb{E}Y_n = \mu n + O(1 + n^{\alpha-1}), \quad m \geq 3, \tag{31}$$

$$333 \quad \text{Var}(Y_n) = \sigma^2 n + O(1 + n^{2\alpha-2}), \quad 3 \leq m \leq 26, \tag{32}$$

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335 Here, the constants  $\mu, \sigma > 0$  depend on  $m$  and  $\alpha \in \mathbb{R}$  depends on  $m$  such that  $\alpha < 1$  for  
 336  $m \leq 13$ ,  $1 \leq \alpha \leq 4/3$  for  $14 \leq m \leq 19$ , and  $4/3 \leq \alpha \leq 3/2$  for  $20 \leq m \leq 26$ , see, e.g.,  
 337 Mahmoud [12, Table 3.1] for the values  $\alpha = \alpha_m$  depending on  $m$ . It is known that  $Y_n$   
 338 standardized by mean and variance satisfies a central limit law for  $m \leq 26$ , whereas the  
 339 standardized sequence has no weak limit for  $m > 26$  due to dominant periodicities, see  
 340 Chern and Hwang [4]. The rate of convergence in the central limit law for  $m \leq 26$  for the  
 341 Kolmogorov metric has been identified in Hwang [9]. Our Theorem 3 implies the central limit  
 342 theorem for  $Y_n$  with  $m \leq 26$  with the same (up to an  $\varepsilon$  for  $3 \leq m \leq 19$ ) rate of convergence  
 343 for the Zolotarev metric  $\zeta_3$ :

344 ► **Theorem 8.** *The size  $Y_n$  of a random  $m$ -ary search tree with  $n$  items inserted satisfies,*  
 345 *for  $m \leq 26$ ,*

$$346 \quad \zeta_3\left(\frac{Y_n - \mathbb{E}Y_n}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = \begin{cases} O(n^{-1/2+\varepsilon}), & 3 \leq m \leq 19, \\ O(n^{-3(3/2-\alpha)}), & 20 \leq m \leq 26, \end{cases} \quad (33)$$

347  
 348 as  $n \rightarrow \infty$ .

349 **Proof.** In order to apply Theorem 3 we have to estimate the orders of  $\|\sum_{r=1}^m (A_r^{(n)})^2 - 1\|_{3/2}$   
 350 and  $\|b^{(n)}\|_3$  with  $A_r^{(n)}$  and  $b^{(n)}$  defined in (3). For this we apply Lemma 7. From (31) and  
 351 (32) we obtain that for the quantities appearing in Lemma 7 we can choose  $f(n) = \mu n$ ,  
 352  $e(n) = 1 \vee n^{\alpha-1}$ ,  $g(n) = \sigma^2 n$ , and  $h(n) = 1 \vee n^{2(\alpha-1)}$ . Hence we obtain

$$353 \quad \left\| \sum_{r=1}^m (\bar{A}_r^{(n)})^2 - 1 \right\|_{3/2} = \left\| \sum_{r=1}^m \frac{I_r^{(n)}}{n} - 1 \right\|_{3/2} = \frac{m-1}{n} = O(n^{-1})$$

354 and  $O(h(n)/g(n)) = O(n^{-(1 \wedge (3-2\alpha))})$ . This implies

$$355 \quad \left\| \sum_{r=1}^m (A_r^{(n)})^2 - 1 \right\|_{3/2}^{3/2} = O(n^{-((3/2) \wedge (3(3/2-\alpha)))}).$$

356 Similarly we obtain

$$357 \quad \|\bar{b}^{(n)}\|_3 = \frac{1}{\sigma\sqrt{n}} \left\| 1 - \mu n + \sum_{r=1}^m \mu I_r^{(n)} \right\|_3 = \frac{1}{\sigma\sqrt{n}} \|1 - \mu(m-1)\|_3 = O(n^{-1/2})$$

358 and  $O(e(n)/g^{1/2}(n)) = O(n^{-(1 \wedge (3/2-\alpha))})$ . This implies

$$359 \quad \|b^{(n)}\|_3^3 = O(n^{-((3/2) \wedge (3(3/2-\alpha)))}).$$

360 Hence, condition (15) is satisfied with  $R(n) = n^{-((3/2) \wedge (3(3/2-\alpha)))}$ . ◀

361 ► **Remark 9.** Using Theorem 1 instead of Theorem 3 in the latter proof is also possible but  
 362 leads to a bound  $O(n^{-(3/2-\alpha)})$  for  $20 \leq m \leq 26$ , missing the factor 3 appearing in Theorem  
 363 8.

364 In the full paper version we also discuss rates of convergence for the number of leaves of  
 365  $d$ -dimensional random point quadrees in the model of [7, 3, 8] where a similar behavior  
 366 as in Theorem 8 appears. A technically related example is the number of maxima in right  
 367 triangles in the model of [1, 2], where the order  $n^{-1/4}$  appears. Our framework also applies.

### 4.3 Periodic functions in mean and variance

We now discuss some examples where the asymptotic expansions of the mean and the variance include periodic functions instead of fixed constants. This is the case for several quantities in binomial splitting processes such as tries, PATRICIA tries and digital search trees. Throughout this section, we assume that we have a 3-integrable sequence  $(Y_n)_{n \geq 0}$  satisfying the recursion

$$Y_n \stackrel{d}{=} Y_{I_1^{(n)}}^{(1)} + Y_{I_2^{(n)}}^{(2)} + b_n, \quad n \geq n_0, \quad (34)$$

with  $(I^{(n)}, b_n)$ ,  $(Y_n^{(1)})_{n \geq 0}$  and  $(Y_n^{(2)})_{n \geq 0}$  independent and  $(Y_n^{(r)})_{n \geq 0} \stackrel{d}{=} (Y_n)_{n \geq 0}$  for  $r = 1, 2$ . Furthermore,  $I_1^{(n)}$  has the binomial distribution  $\text{Bin}(n, \frac{1}{2})$  and  $I_2^{(n)} = n - I_1^{(n)}$  or  $I_1^{(n)}$  is binomially  $\text{Bin}(n - 1, \frac{1}{2})$  distributed and  $I_2^{(n)} = n - 1 - I_1^{(n)}$ . Mostly, these binomial recurrences are asymptotically normally distributed, see [10, 11, 14, 18] for some examples.

Our first theorem covers the case of linear mean and variance, i.e. we assume that, as  $n \rightarrow \infty$ ,

$$\mathbb{E}[Y_n] = nP_1(\log_2 n) + O(1), \quad (35)$$

$$\text{Var}(Y_n) = nP_2(\log_2 n) + O(1), \quad (36)$$

for some smooth and 1-periodic functions  $P_1, P_2$  with  $P_2 > 0$ . Possible applications would start with the analysis of the number of internal nodes of a trie for  $n$  strings in the symmetric Bernoulli model and the number of leaves in a random digital search tree, see, e.g., [10].

► **Theorem 10.** *Let  $(Y_n)_{n \geq 0}$  be 3-integrable and satisfy (34) with  $\|b_n\|_3 = O(1)$ , (35) and (36). Then, for any  $\varepsilon > 0$  and  $n \rightarrow \infty$ , we have*

$$\zeta_3\left(\frac{Y_n - \mathbb{E}[Y_n]}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = O(n^{-1/2+\varepsilon}).$$

We now consider the case where our quantities  $Y_n$  satisfy recursion (34) with  $b_n$  being essentially  $n$ . We assume that, as  $n \rightarrow \infty$ , we have

$$\mathbb{E}[Y_n] = n \log_2(n) + nP_1(\log_2 n) + O(1), \quad (37)$$

$$\text{Var}(Y_n) = nP_2(\log_2 n) + O(1), \quad (38)$$

for some smooth and 1-periodic functions  $P_1, P_2$  with  $P_2 > 0$ . This covers, for example, the external path length of random tries and related digital tree structures constructed from  $n$  random binary strings under appropriate independence assumptions.

► **Theorem 11.** *Let  $(Y_n)_{n \geq 0}$  be 3-integrable and satisfy (34) with  $\|b_n - n\|_3 = O(1)$ , (37) and (38). Then, for any  $\varepsilon > 0$  and  $n \rightarrow \infty$ , we have*

$$\zeta_3\left(\frac{Y_n - \mathbb{E}[Y_n]}{\sqrt{\text{Var}(Y_n)}}, \mathcal{N}(0, 1)\right) = O(n^{-1/2+\varepsilon}).$$

### 4.4 A multivariate application

We consider a random binary search tree with  $n$  nodes built from a random permutation of  $\{1, \dots, n\}$ . For  $n \geq 0$ , we denote by  $L_{0n}$  the number of nodes with no left descendant and

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405 by  $L_{1n}$  the number of nodes with exactly one left descendant. Defining  $Y_n := (L_{0n}, L_{1n})$ , we  
 406 have  $Y_0 = (0, 0)$  and we obtain the following distributional recurrence:

$$407 \quad Y_n \stackrel{d}{=} Y_{I_1^{(n)}}^{(1)} + Y_{I_2^{(n)}}^{(2)} + b_n, \quad n \geq 1,$$

408 where  $(Y_j^{(1)})_{j \geq 0}$  and  $(Y_j^{(2)})_{j \geq 0}$  are independent copies of  $(Y_j)_{j \geq 0}$ ,  $I_1^{(n)}$  is uniformly distributed  
 409 on  $\{0, \dots, n-1\}$  and independent of  $(Y^{(1)})$  and  $(Y^{(2)})$ ,  $I_2^{(n)} = n-1 - I_1^{(n)}$  and  $b_n =$   
 410  $(\mathbf{1}_{\{I_1^{(n)}=0\}}, \mathbf{1}_{\{I_1^{(n)}=1\}})$ . In Devroye [5] it is shown that, for  $n \geq 2$ ,

$$411 \quad \mathbb{E}[L_{0n}] = \frac{1}{2}(n+1), \quad \mathbb{E}[L_{1n}] = \frac{1}{6}(n+1),$$

413 and that the standardized quantities have a limiting normal distribution. Using Devroye's  
 414 description with local counters one also obtains the covariance structure:

415 ► **Lemma 12.** *For  $n \geq 4$ , we have  $\text{Cov}(Y_n) = (n+1)\Gamma$  with*

$$416 \quad \Gamma = \frac{1}{360} \begin{pmatrix} 30 & -15 \\ -15 & 28 \end{pmatrix}.$$

417 For  $n \geq 0$ , we now set  $M_n := \mathbb{E}[Y_n]$ ,  $C_n = \text{Id}_2$  for  $n \leq 3$ ,  $C_n := \text{Cov}(Y_n)$  for  $n \geq 4$  and  
 418 define  $X_n := C_n^{-1/2}(Y_n - M_n)$  for  $n \geq 0$ . Note that the matrix  $\Gamma$  in Lemma 12 is symmetric  
 419 and positive definite, which implies, for  $n \geq 4$ ,

$$420 \quad C_n^{1/2} = \sqrt{n+1} \Gamma^{1/2} \quad \text{and} \quad C_n^{-1/2} = \frac{1}{\sqrt{n+1}} \Gamma^{-1/2}.$$

421 The normalized quantities satisfy  $X_0 = (0, 0)$  and recursion (2) with  $K = 2$ ,  $n_0 = 1$ ,

$$422 \quad A_r^{(n)} = C_n^{-1/2} C_{I_r^{(n)}}^{1/2} = \mathbf{1}_{\{I_r^{(n)} \geq 4\}} \sqrt{\frac{I_r^{(n)} + 1}{n+1}} \text{Id}_2 + \mathbf{1}_{\{I_r^{(n)} < 4\}} \frac{1}{\sqrt{n+1}} \Gamma^{-1/2}$$

423 for  $r = 1, 2$  and

$$424 \quad b^{(n)} = C_n^{-1/2}(b_n - M_n + M_{I_1^{(n)}} + M_{I_2^{(n)}}).$$

425 Modeling all quantities on a joint probability space such that  $I_1^{(n)}/n$  converges almost surely  
 426 to a uniform random variable  $U$  in  $[0, 1]$ , we have the  $L_3$ -convergences  $A_1^{(n)} \rightarrow \sqrt{U} \text{Id}_2$ ,  
 427  $A_2^{(n)} \rightarrow \sqrt{1-U} \text{Id}_2$  and  $b^{(n)} \rightarrow 0$  as  $n \rightarrow \infty$ . Thus, we are in the situation of Section 2.2  
 428 and obtain the limiting equation

$$429 \quad X \stackrel{d}{=} \sqrt{U} X^{(1)} + \sqrt{1-U} X^{(2)},$$

430 with  $U$  uniformly distributed on  $[0, 1]$  and  $X^{(1)}$ ,  $X^{(2)}$  and  $U$  independent. We now check the  
 431 conditions of Theorem 3. Since  $A_1^{(n)}(A_1^{(n)})^T + A_2^{(n)}(A_2^{(n)})^T = \text{Id}_2$  on the event  $\{I_1^{(n)}, I_2^{(n)} \geq 4\}$ ,  
 432 we obtain, as  $n \rightarrow \infty$ ,

$$433 \quad \left\| \sum_{r=1}^2 A_r^{(n)}(A_r^{(n)})^T - \text{Id}_2 \right\|_{3/2}^{3/2} = O\left( \left\| \mathbf{1}_{\{I_1^{(n)} < 4\}} \left( \frac{1}{n+1} \Gamma^{-1} + \frac{I_2^{(n)} + 1}{n+1} \text{Id}_2 - \text{Id}_2 \right) \right\|_{3/2}^{3/2} \right)$$

$$434 \quad = O\left( \mathbb{E} \left[ \mathbf{1}_{\{I_1^{(n)} < 4\}} \left\| \frac{1}{n+1} \Gamma^{-1} - \frac{I_1^{(n)} + 1}{n+1} \text{Id}_2 \right\|_{\text{op}}^{3/2} \right] \right)$$

$$435 \quad = O(n^{-5/2}).$$

437 Similarly, we obtain

$$438 \quad \|b^{(n)}\|_3^3 = O(n^{-5/2}).$$

440 Since we have  $\|\mathbf{1}_{\{I_r^{(n)} < \ell\}} A_r^{(n)}\|_3^3 = O(n^{-5/2})$  for  $\ell \in \mathbb{N}$  and  $r = 1, 2$ , the technical conditions  
 441 are satisfied. We now use Theorem 3 with  $R(n) = n^{-1/2}$ . Note that condition (16) is not  
 442 satisfied for  $R(n) = n^{-1/2}$ , but we can use the weakened condition stated in Remark 2 to  
 443 obtain the following result.

444 ► **Theorem 13.** Denoting by  $Y_n := (L_{0n}, L_{1n})$  the vector of the numbers of nodes with no  
 445 and with exactly one left descendant respectively in a random binary search tree with  $n$  nodes  
 446 we have, for  $n \rightarrow \infty$ , that

$$447 \quad \zeta_3(\text{Cov}(Y_n)^{-1/2}(Y_n - \mathbb{E}[Y_n]), \mathcal{N}(0, \text{Id}_2)) = O(n^{-1/2}).$$

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## 5 Appendix

508 **Proof.** (*Proof of Theorem 1*) Using condition (10), the assumption that  $R$  is monotonically  
509 decreasing and condition (11), we have

$$510 \quad \mathbb{E} \left[ \sum_{r=1}^K \|A_r^*\|_{\text{op}}^s \right] = \lim_{n \rightarrow \infty} \mathbb{E} \left[ \sum_{r=1}^K \|A_r^{(n)}\|_{\text{op}}^s \right] \leq \limsup_{n \rightarrow \infty} \mathbb{E} \left[ \sum_{r=1}^K \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right] < 1.$$

512 Furthermore, condition (10) implies  $\mathbb{E}[b^*] = \lim_{n \rightarrow \infty} \mathbb{E}[b^{(n)}] = 0$  if  $s > 1$  and additionally

$$513 \quad \mathbb{E}[b^*(b^*)^T] + \mathbb{E} \left[ \sum_{r=1}^K A_r^*(A_r^*)^T \right] = \text{Id}_d$$

514 if  $s > 2$ . Thus, Corollary 3.4 in [14] states that equation (12) has a unique fixed-point  
515  $\mathcal{L}(X)$  in  $\mathcal{P}_s^d(0, \text{Id}_d)$ . To establish a rate of convergence to this fixed-point, we introduce the  
516 accompanying sequence

$$517 \quad Z_n^* := \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} X^{(r)} + b^{(n)},$$

518 where  $(A_1^{(n)}, \dots, A_K^{(n)}, I^{(n)}, b^{(n)}, X^{(1)}, \dots, X^{(K)})$  are independent and  $X^{(r)}$  is identically dis-  
519 tributed as  $X$  for  $r = 1, \dots, K$ . Here, for  $2 < s \leq 3$ , the sequence  $(T_n)_{n \geq 0}$  is chosen such  
520 that  $Z_n^*$  has the same covariance structure as  $X_n$ . To be more precise, for  $2 < s \leq 3$ , we  
521 choose  $T_n$  such that  $T_n T_n^T = \text{Cov}(X_n)$  (i.e.  $T_n = \text{Id}_d$  for  $n \geq n_1$  and  $T_n T_n^T = \text{Cov}(Y_n)$  for  
522  $n < n_1$ ). For  $s \leq 2$ , we do not need to control the covariance of  $Z_n^*$  and set  $T_n := \text{Id}_d$   
523 for  $n \geq 0$ . Then,  $Z_n^*$  is  $L_s$ -integrable, we have  $\mathbb{E}[Z_n^*] = 0$  for  $s > 1$  and in the case  $s > 2$

524 additionally  $\text{Cov}(Z_n^*) = \text{Cov}(X_n) = \text{Id}_d$  for  $n \geq n_1$ . Hence,  $\zeta_s$ -distances between  $X_n$ ,  $Z_n^*$   
525 and  $X$  are finite for  $n \geq n_1$ . Applying the triangle inequality we have, for  $n \geq n_1$ ,

$$526 \quad \zeta_s(X_n, X) \leq \zeta_s(X_n, Z_n^*) + \zeta_s(Z_n^*, X). \quad (39)$$

528 Denoting by  $\Upsilon_n$  the joint distribution of  $(A_1^{(n)}, \dots, A_K^{(n)}, b^{(n)}, I^{(n)})$ ,  $\alpha = (\alpha_1, \dots, \alpha_K)$ ,  $j =$   
529  $(j_1, \dots, j_K)$  and  $\Delta(n) := \zeta_s(X_n, X)$ , we obtain by conditioning on  $\Upsilon_n$  that, for  $n \geq n_1$ ,

$$530 \quad \begin{aligned} \zeta_s(X_n, Z_n^*) &= \zeta_s \left( \sum_{r=1}^K A_r^{(n)} X_{I_r^{(n)}}^{(r)} + b^{(n)}, \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} X^{(r)} + b^{(n)} \right) \\ 531 \quad &= \sup_{f \in \mathcal{F}_s} \left| \int \mathbb{E} \left[ f \left( \sum_{r=1}^K \alpha_r X_{j_r}^{(r)} + \beta \right) \right] - \mathbb{E} \left[ f \left( \sum_{r=1}^K \alpha_r T_{j_r} X^{(r)} + \beta \right) \right] d\Upsilon_n(\alpha, \beta, j) \right| \\ 532 \quad &\leq \int \zeta_s \left( \sum_{r=1}^K \alpha_r X_{j_r}^{(r)} + \beta, \sum_{r=1}^K \alpha_r T_{j_r} X^{(r)} + \beta \right) d\Upsilon_n(\alpha, \beta, j) \\ 533 \quad &\leq \int \sum_{r=1}^K \|\alpha_r\|_{\text{op}}^s \zeta_s \left( X_{j_r}^{(r)}, T_{j_r} X^{(r)} \right) d\Upsilon_n(\alpha, \beta, j) \\ 534 \quad &\leq \left( \mathbb{E} \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)}=n\}} \|A_r^{(n)}\|_{\text{op}}^s \right) \Delta(n) + \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{n_1 \leq I_r^{(n)} < n\}} \|A_r^{(n)}\|_{\text{op}}^s \Delta(I_r^{(n)}) \right] \\ 535 \quad &\quad + \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)} < n_1\}} \|A_r^{(n)}\|_{\text{op}}^s \sup_{k < n_1} \zeta_s(X_k, T_k X^{(r)}) \right]. \quad (40) \\ 536 \end{aligned}$$

537 Note that the last summand is in  $O(R(n))$  by condition (8). To bound the second summand  
538  $\zeta_s(Z_n^*, X)$  in (39), we switch to the Wasserstein metric  $\ell_s$ : By condition (10) and  $\|Z_n^*\|_s \leq$   
539  $\sum_{r=1}^K \|A_r^{(n)} T_{I_r^{(n)}}\|_s \|X\|_s + \|b^{(n)}\|_s$ , we have  $\sup_{n \geq 0} \|Z_n^*\|_s < \infty$ . Thus, a standard bound  
540 implies that  $\zeta_s(Z_n^*, X) \leq C_s \ell_s(Z_n^*, X)$  for some constant  $C_s > 0$ . Furthermore, we have

$$541 \quad \begin{aligned} \ell_s(Z_n^*, X) &\leq \left\| \left( \sum_{r=1}^K A_r^{(n)} T_{I_r^{(n)}} X^{(r)} + b^{(n)} \right) - \left( \sum_{r=1}^K A_r^* X^{(r)} + b^* \right) \right\|_s \\ 542 \quad &\leq \sum_{r=1}^K \|A_r^{(n)} T_{I_r^{(n)}} - A_r^*\|_s \|X^{(r)}\|_s + \|b^{(n)} - b^*\|_s \\ 543 \quad &\leq \sum_{r=1}^K \left( \|A_r^{(n)} T_{I_r^{(n)}} - A_r^{(n)}\|_s + \|A_r^{(n)} - A_r^*\|_s \right) \|X\|_s + \|b^{(n)} - b^*\|_s \\ 544 \quad &= \sum_{r=1}^K \left( \|\mathbf{1}_{\{I_r^{(n)} < n_1\}} A_r^{(n)} (T_{I_r^{(n)}} - \text{Id}_d)\|_s + \|A_r^{(n)} - A_r^*\|_s \right) \|X\|_s + \|b^{(n)} - b^*\|_s. \\ 545 \end{aligned}$$

546 Using conditions (8) and (10), we obtain  $\ell_s(Z_n^*, X) = O(R(n))$ . Hence, putting everything  
547 together and introducing the notation  $p_n := \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{I_r^{(n)}=n\}} \|A_r^{(n)}\|_{\text{op}}^s \right]$ , we obtain from  
548 (39) and (40) that

$$549 \quad \Delta(n) \leq p_n \Delta(n) + \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{n_1 \leq I_r^{(n)} < n\}} \|A_r^{(n)}\|_{\text{op}}^s \Delta(I_r^{(n)}) \right] + O(R(n)). \quad (41)$$

551 From (11), there exists a  $\delta > 0$  such that  $\mathbb{E} \left[ \sum_{r=1}^K \frac{R(I_r^{(n)})}{R(n)} \|A_r^{(n)}\|_{\text{op}}^s \right] \leq 1 - \delta$  for all  $n$   
552 sufficiently large and from (9) we have  $p_n < \delta/2$  for  $n$  large. We now choose some  $C > 0$

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553 and  $n_2 \geq n_1$  sufficiently large such that for  $n \geq n_2$  all these inequalities are satisfied and the  
 554  $O(R(n))$  term in (41) is bounded by  $CR(n)$ . By setting

$$555 \quad L := \frac{2C}{\delta} \vee \max \left\{ \frac{\Delta(n)}{R(n)} : n \leq n_2 \right\}$$

556 we now obtain  $\Delta(n) \leq LR(n)$  by induction: For  $n \leq n_2$ , by definition of  $L$ , the assertion is  
 557 true. For  $n > n_2$ , solving for  $\Delta(n)$  in (41), we find

$$\begin{aligned} 558 \quad \Delta(n) &\leq \frac{1}{1-p_n} \left( \mathbb{E} \left[ \sum_{r=1}^K \mathbf{1}_{\{n_1 \leq I_r^{(n)} < n\}} \|A_r^{(n)}\|_{\text{op}}^s \Delta(I_r^{(n)}) \right] + CR(n) \right) \\ 559 \quad &\leq \frac{1}{1-\delta/2} \left( \mathbb{E} \left[ \sum_{r=1}^K \|A_r^{(n)}\|_{\text{op}}^s LR(I_r^{(n)}) \right] + CR(n) \right) \\ 560 \quad &= \frac{1}{1-\delta/2} \left( L \mathbb{E} \left[ \sum_{r=1}^K \|A_r^{(n)}\|_{\text{op}}^s \frac{R(I_r^{(n)})}{R(n)} \right] R(n) + CR(n) \right) \\ 561 \quad &\leq \frac{1}{1-\delta/2} (L(1-\delta) + C) R(n) \\ 562 \quad &\leq LR(n). \end{aligned}$$

564

565 **Proof.** (*Proof of Lemma 5*) As the matrix  $G_n G_n^T$  is symmetric and positive-semidefinite, we  
 566 can decompose it in the following way: Let  $\lambda_1 \geq \dots \geq \lambda_m \geq 1 > \lambda_{m+1} \geq \dots \geq \lambda_d \geq 0$  be  
 567 the (random) eigenvalues of  $G_n G_n^T$ . Then, with a suitable (random) orthogonal matrix  $O$ ,  
 568 we have

$$\begin{aligned} 569 \quad G_n G_n^T &= O \text{diag}(\lambda_1, \dots, \lambda_d) O^T \\ 570 \quad &= O \text{diag}(1, \dots, 1, \lambda_{m+1}, \dots, \lambda_d) O^T + O \text{diag}(\lambda_1 - 1, \dots, \lambda_m - 1, 0, \dots, 0) O^T \\ 571 \quad &= B_n B_n^T + C_n C_n^T, \end{aligned}$$

573 where we define the random  $(d \times d)$ -matrices  $B_n := O \text{diag}(1, \dots, 1, \sqrt{\lambda_{m+1}}, \dots, \sqrt{\lambda_d}) O^T$   
 574 and  $C_n := O \text{diag}(\sqrt{\lambda_1 - 1}, \dots, \sqrt{\lambda_m - 1}, 0, \dots, 0) O^T$ . Hence, we can decompose  $Z_n^*$  in the  
 575 following way:

$$576 \quad Z_n^* \stackrel{d}{=} G_n N + b^{(n)} \stackrel{d}{=} B_n N + C_n N' + b^{(n)} =: \hat{Z}_n^*,$$

577 where  $(B_n, C_n, b^{(n)})$ ,  $N$  and  $N'$  are independent with  $\mathcal{L}(N) = \mathcal{L}(N') = \mathcal{N}(0, \text{Id}_d)$ . Analog-  
 578 ously, we decompose the multivariate normal distribution:

$$579 \quad N \stackrel{d}{=} B_n N + D_n N' =: \hat{N},$$

580 where  $D_n := O \text{diag}(0, \dots, 0, \sqrt{1 - \lambda_{m+1}}, \dots, \sqrt{1 - \lambda_d}) O^T$  is chosen such that  $B_n B_n^T +$   
 581  $D_n D_n^T = \text{Id}_d$ .

582 By definition of the Zolotarev metric  $\zeta_3$  we have

$$583 \quad \zeta_3(Z_n^*, \mathcal{N}(0, \text{Id}_d)) = \zeta_3(\hat{Z}_n^*, \hat{N}) = \sup_{f \in \mathcal{F}_3} \left| \mathbb{E}[f(\hat{Z}_n^*) - f(\hat{N})] \right|.$$

584 For arbitrary  $f \in \mathcal{F}_3$  we use Taylor expansion around  $N$  and obtain for  $x \in \mathbb{R}^d$  that

$$585 \quad f(x) = f(N) + (x - N)^T \nabla f(N) + \frac{1}{2} (x - N)^T H_f(N) (x - N) + R(x, N),$$



586 where the remainder term satisfies  $|R(x, N)| \leq \frac{1}{2} \|x - N\|^3$ . Thus, we have

$$587 \quad f(\hat{Z}_n^*) - f(\hat{N}) = (\hat{Z}_n^* - \hat{N})^T \nabla f(N) + \frac{1}{2} (\hat{Z}_n^* - N)^T H_f(N) (\hat{Z}_n^* - N) \\ 588 \quad \quad \quad - \frac{1}{2} (\hat{N} - N)^T H_f(N) (\hat{N} - N) + R(\hat{Z}_n^*, N) - R(\hat{N}, N). \quad (42) \\ 589$$

590 We now study the expectation of these summands: For the first summand, we have

$$591 \quad \mathbb{E}[(\hat{Z}_n^* - \hat{N})^T \nabla f(N)] = \mathbb{E}[(C_n - D_n)N' + b^{(n)}]^T \nabla f(N) \\ 592 \quad \quad \quad = \mathbb{E}[(C_n - D_n)N' + b^{(n)}]^T \mathbb{E}[\nabla f(N)] = 0,$$

594 since  $N$  is independent of the other quantities,  $N'$  is independent of  $(C_n, D_n)$  and  $\mathbb{E}[N'] =$   
595  $\mathbb{E}[b^{(n)}] = 0$ . For the second summand, we define  $F_n := B_n - \text{Id}_d$  and obtain

$$596 \quad \mathbb{E}[(\hat{Z}_n^* - N)^T H_f(N) (\hat{Z}_n^* - N)] \\ 597 \quad = \mathbb{E}[(F_n N + C_n N' + b^{(n)})^T H_f(N) (F_n N + C_n N' + b^{(n)})] \\ 598 \quad = \mathbb{E}[(F_n N)^T H_f(N) (F_n N)] + \mathbb{E}[(F_n N)^T H_f(N) (C_n N')] + \mathbb{E}[(F_n N)^T H_f(N) b^{(n)}] \\ 599 \quad \quad + \mathbb{E}[(C_n N')^T H_f(N) (F_n N)] + \mathbb{E}[(C_n N')^T H_f(N) (C_n N')] + \mathbb{E}[(C_n N')^T H_f(N) b^{(n)}] \\ 600 \quad \quad + \mathbb{E}[(b^{(n)})^T H_f(N) (F_n N)] + \mathbb{E}[(b^{(n)})^T H_f(N) (C_n N')] + \mathbb{E}[(b^{(n)})^T H_f(N) b^{(n)}]. \\ 601$$

602 Since  $N, N'$  and  $(F_n, C_n, b^{(n)})$  are independent with  $\mathbb{E}[N'] = 0$ , we have

$$603 \quad \mathbb{E}[(F_n N)^T H_f(N) (C_n N')] = 0.$$

604 The same argument applies to  $\mathbb{E}[(C_n N')^T H_f(N) (F_n N)]$ ,  $\mathbb{E}[(C_n N')^T H_f(N) b^{(n)}]$  and  $\mathbb{E}[(b^{(n)})^T H_f(N) (C_n N')]$ .  
605 Analogously, we obtain for the third summand in (42)

$$606 \quad \mathbb{E}[(\hat{N} - N)^T H_f(N) (\hat{N} - N)] \\ 607 \quad = \mathbb{E}[(F_n N + D_n N')^T H_f(N) (F_n N + D_n N')] \\ 608 \quad \quad = \mathbb{E}[(F_n N)^T H_f(N) (F_n N)] + \mathbb{E}[(D_n N')^T H_f(N) (D_n N')]. \\ 609$$

610 This implies together with  $\mathbb{E}[(F_n N)^T H_f(N) b^{(n)}] = \mathbb{E}[(b^{(n)})^T H_f(N) (F_n N)]$

$$611 \quad \mathbb{E}[(\hat{Z}_n^* - N)^T H_f(N) (\hat{Z}_n^* - N) - (\hat{N} - N)^T H_f(N) (\hat{N} - N)] \\ 612 \quad = \mathbb{E}[(C_n N')^T H_f(N) (C_n N')] - \mathbb{E}[(D_n N')^T H_f(N) (D_n N')] + \mathbb{E}[(b^{(n)})^T H_f(N) b^{(n)}] \\ 613 \quad \quad + 2 \mathbb{E}[(F_n N)^T H_f(N) b^{(n)}] \\ 614$$

615 Note that we have  $C_n C_n^T - D_n D_n^T = G_n G_n^T - \text{Id}_d$ . Furthermore,  $\mathbb{E}[G_n G_n^T + b^{(n)} (b^{(n)})^T] =$   
616  $\text{Id}_d$ . Thus, with the independence of  $N, N'$  and  $(C_n, D_n, b^{(n)})$  and  $\mathbb{E}[N'_i N'_j] = \mathbf{1}_{\{i=j\}}$  for  
617  $i, j = 1, \dots, d$ , we have

$$618 \quad \mathbb{E}[(C_n N')^T H_f(N) (C_n N')] - \mathbb{E}[(D_n N')^T H_f(N) (D_n N')] + \mathbb{E}[(b^{(n)})^T H_f(N) b^{(n)}] \\ 619 \quad = \sum_{i,j=1}^d \mathbb{E}[H_f(N)_{ij}] \mathbb{E}[(C_n N')_i (C_n N')_j - (D_n N')_i (D_n N')_j + b_i^{(n)} b_j^{(n)}] \\ 620 \quad = \sum_{i,j=1}^d \mathbb{E}[H_f(N)_{ij}] \mathbb{E}[(C_n C_n^T - D_n D_n^T)_{ij} + (b^{(n)} (b^{(n)})^T)_{ij}] \\ 621 \quad = \sum_{i,j=1}^d \mathbb{E}[H_f(N)_{ij}] \mathbb{E}[(G_n G_n^T + b^{(n)} (b^{(n)})^T - \text{Id}_d)_{ij}] \\ 622 \quad \quad = 0. \\ 623$$

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624 Thus, we have shown that

$$\begin{aligned}
 625 \quad \left| \mathbb{E}[f(\hat{Z}_n^*) - f(\hat{N})] \right| &= \left| \mathbb{E}[(F_n N)^T H_f(N) b^{(n)}] + \mathbb{E}[R(\hat{Z}_n^*, N)] - \mathbb{E}[R(\hat{N}, N)] \right| \\
 626 \quad &\leq \mathbb{E}[|(F_n N)^T H_f(N) b^{(n)}|] + \mathbb{E}[|R(\hat{Z}_n^*, N)|] + \mathbb{E}[|R(\hat{N}, N)|].
 \end{aligned}$$

628 We now bound these three terms. For this, without loss of generality, we may assume that  
 629  $H_f(0) = 0$ : If this is not the case, we consider the function  $g : \mathbb{R}^d \rightarrow \mathbb{R}$  defined by  $g(x) :=$   
 630  $f(x) - \frac{1}{2} x^T H_f(0) x$  for  $x \in \mathbb{R}^d$ . Then,  $H_g(0) = 0$  and  $\mathbb{E}[g(\hat{Z}_n^*) - g(\hat{N})] = \mathbb{E}[f(\hat{Z}_n^*) - f(\hat{N})]$  since  
 631  $\hat{Z}_n^*$  and  $\hat{N}$  have the same mean and covariance structure. The assumption  $H_f(0) = 0$  implies,  
 632 together with the Lipschitz property of the second derivative of  $f$ ,  $\|H_f(N)\|_{\text{op}} \leq \|N\|$ . Hence,  
 633 using the Cauchy-Schwarz inequality, the independence of  $(F_n, b^{(n)})$  and  $N$  and Hölder's  
 634 inequality, we have

$$\begin{aligned}
 635 \quad \mathbb{E}[|(F_n N)^T H_f(N) b^{(n)}|] &\leq \mathbb{E}[\|F_n\|_{\text{op}} \|N\| \|H_f(N)\|_{\text{op}} \|b^{(n)}\|] \\
 636 \quad &\leq \mathbb{E}[\|N\|^2] \mathbb{E}[\|F_n\|_{\text{op}} \|b^{(n)}\|] \\
 637 \quad &\leq d \|F_n\|_{3/2} \|b^{(n)}\|_3 \\
 638 \quad &\leq d \|G_n G_n^T - \text{Id}_d\|_{3/2} \|b^{(n)}\|_3,
 \end{aligned}$$

640 where the last step follows by  $\|G_n G_n^T - \text{Id}_d\|_{\text{op}} = \max\{|\lambda_1 - 1|, |\lambda_d - 1|\}$ ,  $\|F_n\|_{\text{op}} =$   
 641  $\mathbf{1}_{\{\lambda_d < 1\}} |\sqrt{\lambda_d} - 1|$  and the identity  $|\sqrt{a} - 1| \leq |a - 1|$  for  $a \geq 0$ . The first remainder  
 642 term is bounded by

$$\begin{aligned}
 643 \quad \mathbb{E}[|R(\hat{Z}_n^*, N)|] &\leq \frac{1}{2} \mathbb{E}[\|\hat{Z}_n^* - N\|^3] \\
 644 \quad &= \frac{1}{2} \mathbb{E}[\|F_n N + C_n N' + b^{(n)}\|^3] \\
 645 \quad &= \mathcal{O}(\mathbb{E}[\|F_n\|_{\text{op}}^3] + \mathbb{E}[\|C_n\|_{\text{op}}^3] + \mathbb{E}[\|b^{(n)}\|^3]) \\
 646 \quad &= \mathcal{O}(\|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3),
 \end{aligned}$$

648 since  $\|C_n\|_{\text{op}} = \mathbf{1}_{\{\lambda_1 > 1\}} \sqrt{|\lambda_1 - 1|} \leq \|G_n G_n^T - \text{Id}_d\|_{\text{op}}^{1/2}$  and  $\|F_n\|_{\text{op}} = \mathbf{1}_{\{\lambda_d < 1\}} |\sqrt{\lambda_d} - 1| \leq$   
 649  $\|G_n G_n^T - \text{Id}_d\|_{\text{op}}^{1/2}$  (note that we have  $|\sqrt{a} - 1| \leq \sqrt{|a - 1|}$  for any  $a \geq 0$ ). With the same  
 650 arguments, we obtain for the second remainder term

$$\begin{aligned}
 651 \quad \mathbb{E}[|R(\hat{N}, N)|] &\leq \frac{1}{2} \mathbb{E}[\|F_n N + D_n N'\|^3] = \mathcal{O}(\|F_n\|_3^3 + \|D_n\|_3^3) \\
 652 \quad &= \mathcal{O}(\|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2}),
 \end{aligned}$$

654 as  $\|D_n\|_{\text{op}} = \mathbf{1}_{\{\lambda_d < 1\}} \sqrt{|\lambda_d - 1|} \leq \|G_n G_n^T - \text{Id}_d\|_{\text{op}}^{1/2}$ . This implies

$$\begin{aligned}
 655 \quad \left| \mathbb{E}[f(\hat{Z}_n^*) - f(\hat{N})] \right| &\leq \mathbb{E}[|(F_n N)^T H_f(N) b^{(n)}|] + \mathbb{E}[|R(\hat{Z}_n^*, N)|] + \mathbb{E}[|R(\hat{N}, N)|] \\
 656 \quad &= \mathcal{O}(\|G_n G_n^T - \text{Id}_d\|_{3/2} \|b^{(n)}\|_3 + \|G_n G_n^T - \text{Id}_d\|_{3/2}^{3/2} + \|b^{(n)}\|_3^3) \\
 657 \quad &= \mathcal{O}(R(n)).
 \end{aligned}$$

659 Note that the constants in the O-notation do not depend on the function  $f$ , i.e. we have  
 660  $\sup_{f \in \mathcal{F}_3} \left| \mathbb{E}[f(\hat{Z}_n^*) - f(\hat{N})] \right| = \mathcal{O}(R(n))$ . ◀